

Lecture 13 - Oct 23

Bridge Controller

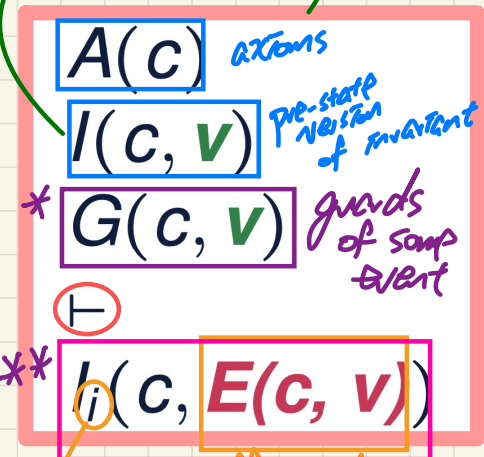
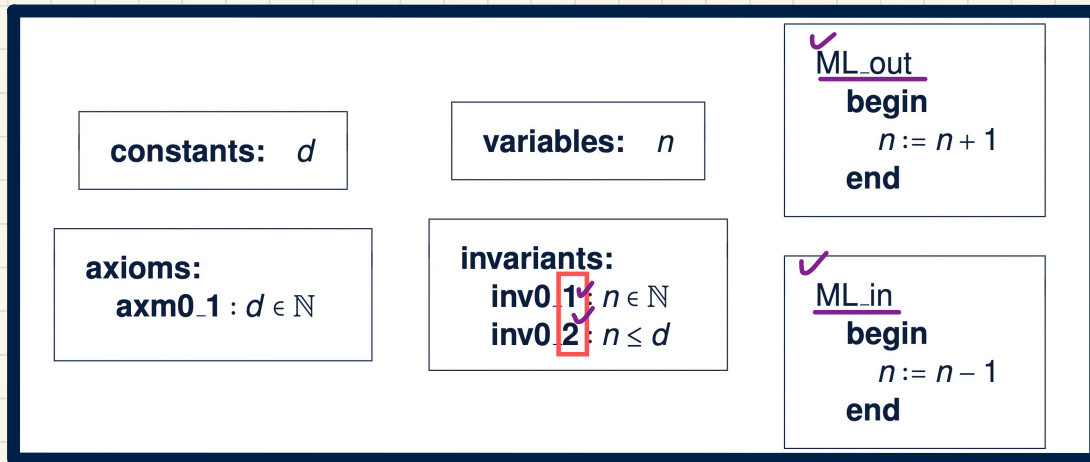
Proof Obligation Rule: Inv. Preservation
Inference Rule: Syntax and Semantics
Sequent Proofs via Inference Rules

Announcements/Reminders

- Today's class: [notes template](#) posted
- **ProgTest** results to be released by next Tuesday's class
- **WrittenTest1** results to be released by early Monday

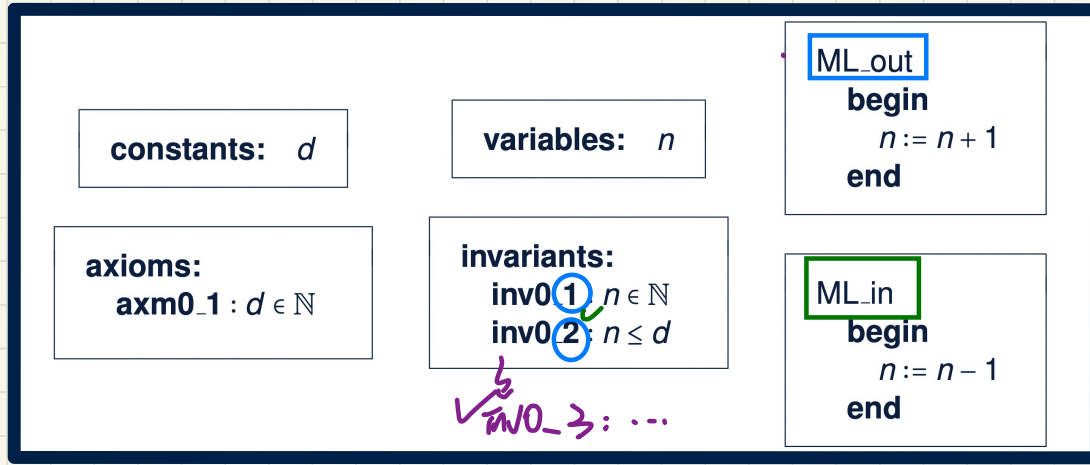
PO/VC Rule of Invariant Preservation: Sequents

assume: all invariants hold in pre-state



Q. How many PO/VC rules for model m0?

- sequents*
- * Guards of some events: ML_out or ML_in (2)
 - ** Some invariant condition: inv0_1 or inv0_2 (2)
- Total # of POs (sequents): $2 \times 2 = 4$



$A(c)$
 $I(c, \mathbf{v})$
 $G(c, \mathbf{v})$
 $\vdash$
 $I(c, \mathbf{E}(c, \mathbf{v}))$

P01: ML_out / inv0_1 / INV invariant preservation

P02: ML_out / inv0_2 / INV

P03: ML_in / inv0_1 / INV

P04: ML_in / inv0_2 / INV

P05: ML_out / inv0_3 / INV

P06: ML_in / inv0_3 / INV

ML_out / inv0-1 / INV

↳ P.O. is related to whether or not taking a state transition via some event action can preserve/maintain inv0-1.

PO/VC Rule of Invariant Preservation: Sequents

constants: d

variables: n

axioms:
 $\text{axm0_1} : d \in \mathbb{N}$

invariants:
 $\text{inv0_1} : n \in \mathbb{N}$
 $\text{inv0_2} : n \leq d$

ML_out $\hookleftarrow$ true
begin
 $n := n + 1$
end
BAP: $n' = n + 1$

ML_in $\hookleftarrow$ true
begin
 $n := n - 1$
end
BAP: $n' = n - 1$

$A(c)$

$I(c, v)$

$G(c, v)$

$\vdash$

$I_i(c, E(c, v))$

Pol: $\text{ML_out} / \text{inv0_1} / \text{INV}$

$d \in \mathbb{N}$
 $n \in \mathbb{N}$
 $n \leq d$
 true
 $\vdash n+1 \in \mathbb{N}$

Poz: $\text{ML_in} / \text{inv0_2} / \text{INV}$

$d \in \mathbb{N}$
 $n \in \mathbb{N}$
 $n \leq d$
 true
 $\vdash n-1 \leq d$

EXERCISES

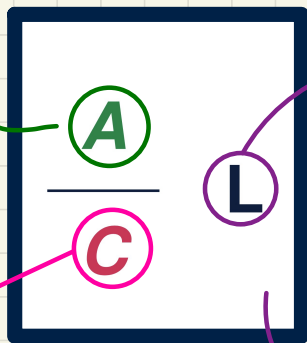
Poz: $\text{ML_out} / \text{inv0_2} / \text{INV}$

Pol: $\text{ML_in} / \text{inv0_1} / \text{INV}$

1. Copy the TN.
2. adapt it to post-step version
3. substitute each primed var. using BAP

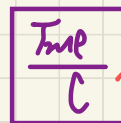
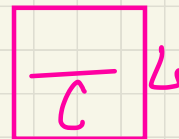
Inference Rule: Syntax and Semantics

Syntax



Semantics

$$A \Rightarrow C \equiv \text{True.}$$

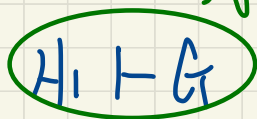


means C is an axiom

Q. What does it mean when A is empty/absent?

without any proof.

Examples



a single sequent.

To prove the consequent C ,

it's sufficient to prove the antecedent A

rewriting **instead**

To prove $H1, H2 \vdash G$ we can drop hypotheses.
it's sufficient to prove $H1 \vdash G$

MON
 $\hookrightarrow$
monotonicity

Sequent

$$\boxed{\begin{array}{c} H \\ \vdash \\ G \end{array}}$$

$$H \Rightarrow G$$

(subject to prove)

e.g.

$$\boxed{\begin{array}{c} H_1 \\ H_2 \\ \vdash \\ G \end{array}}$$

now

$$\boxed{\begin{array}{c} H_1 \\ \vdash \\ G \end{array}}$$

Inference Rule

$$\boxed{\frac{A}{C} \vdash}$$

may be used as a proof step.
↳ to make progress in proofs, rewrite C as A

$$A \Rightarrow C \equiv \text{True}$$

↳ (assumed to be true)

Proof of Sequent: Steps and Structure

Outstanding Sequent to Prove

$d \in \mathbb{N}$
 $n \in \mathbb{N}$
 $n \leq d$
 $\vdash$
 $n + 1 \in \mathbb{N}$

ML_out/inv0_1/INV

Known Inference Rules

$\checkmark$
 $H1 \vdash G$
 $\hline H1, H2 \vdash G$ MON

rewriting

$n \in \mathbb{N} \vdash n + 1 \in \mathbb{N}$ P2

best case

relevant, useful

can be dropped

$d \in \mathbb{N}$
 $n \in \mathbb{N}$ ✓
 $n \leq d$
 $\vdash$
 $n + 1 \in \mathbb{N}$

not relevant

not useful for arguing $n+1 \in \mathbb{N}$

MON

goal

$n \in \mathbb{N}$
 $\vdash$
 $n + 1 \in \mathbb{N}$

P2

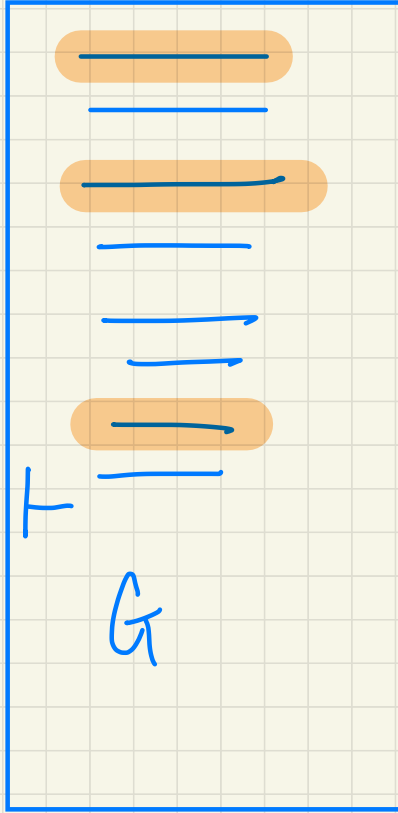
* Analyze the goal predicate to

prop :

Q1. what are the relevant vars?

Q2. what hypotheses are useful?

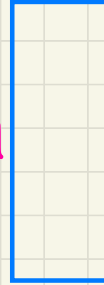
out standing sequent



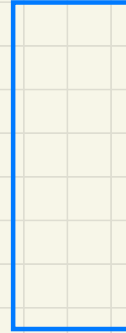
MON



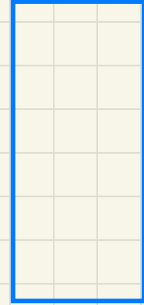
R_1



R_2



...



loop resp



A



axiom rule



Understanding Inference Rule: OR_L

OR_L
OR_R
AND_L
AND_R

$$\frac{H, P \vdash R \quad H, Q \vdash R}{H, P \vee Q \vdash R} \text{OR}_L$$

Q. Does OR_L help us: $\text{AND} \vdash R$

(A) split one sequent to prove into two sequents to prove

(B) combine two sequents to prove into one sequent to prove.

$$\frac{A}{C}$$

To prove C, sufficient to prove A

disjunction

OR appears to the L of $\vdash$

